

IP ROUTING

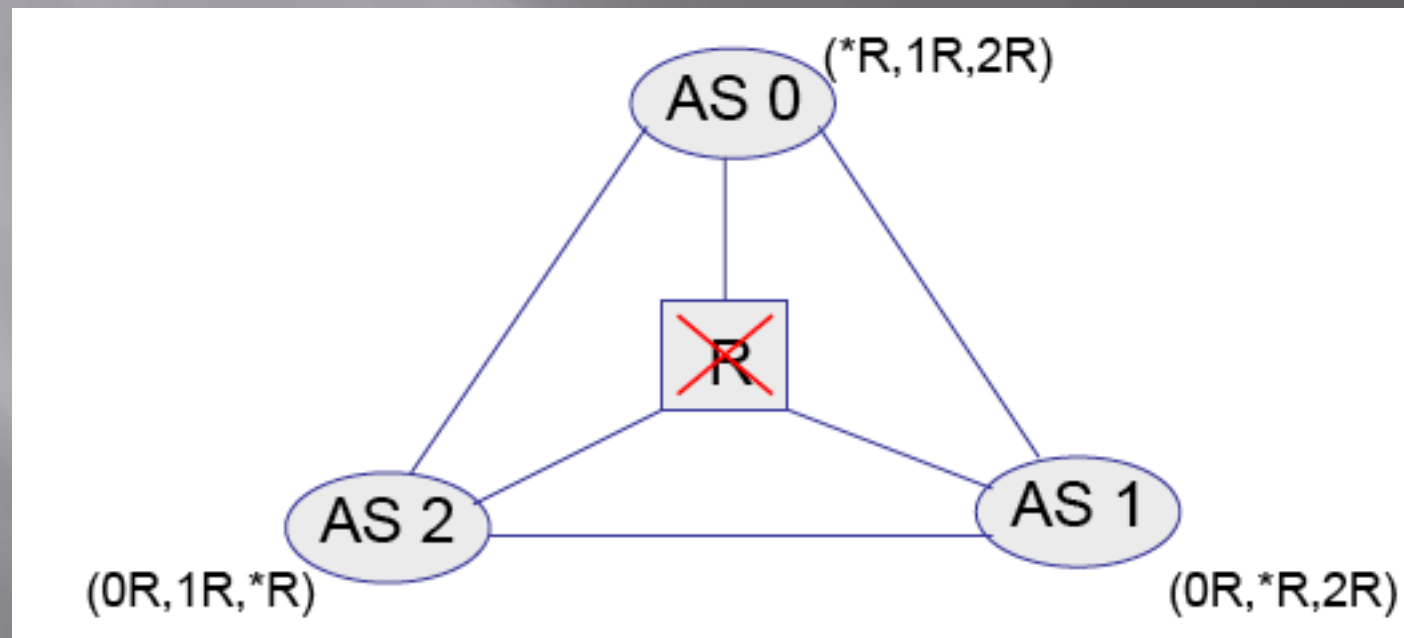
Observations

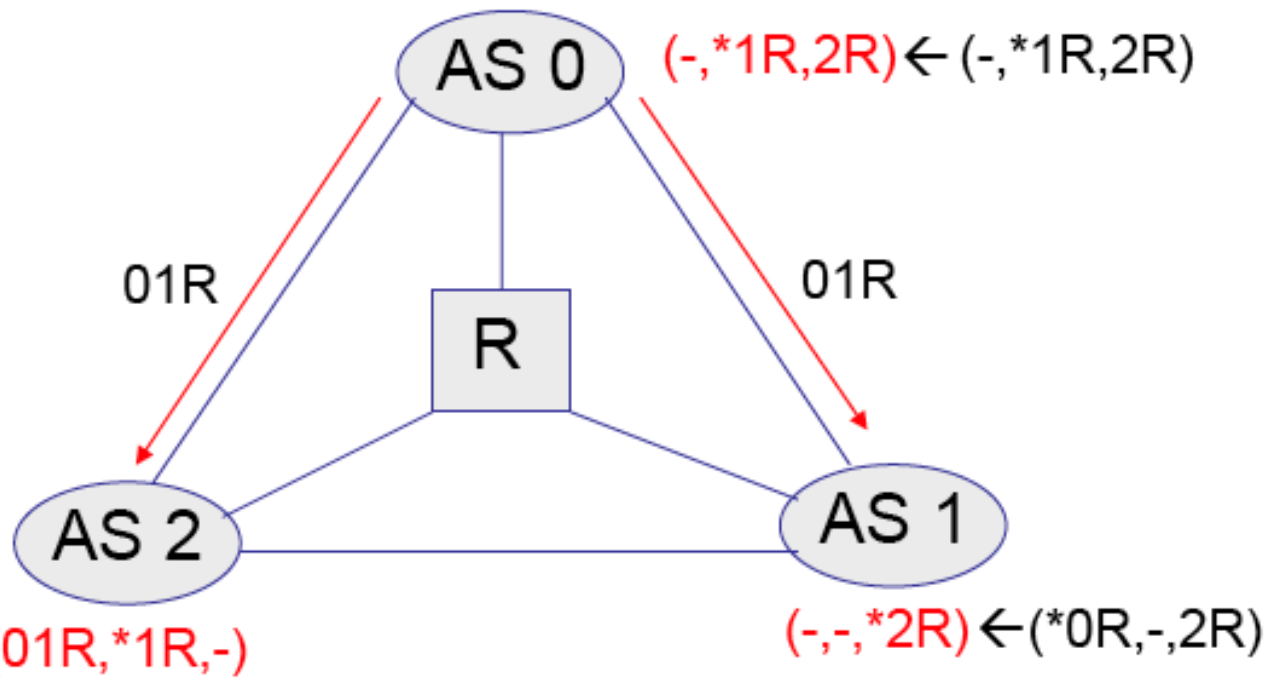
- ▣ Repairs (Tup time) exhibit similar convergence properties as long-short ASPath fail-over
- ▣ Failures (Tdown) and short-long fail-overs (e.g. primary to secondary paths) also similar
 - Slower than a repair (bad news doesn't travel fast, DV protocols)
 - 60% take > 2 minutes

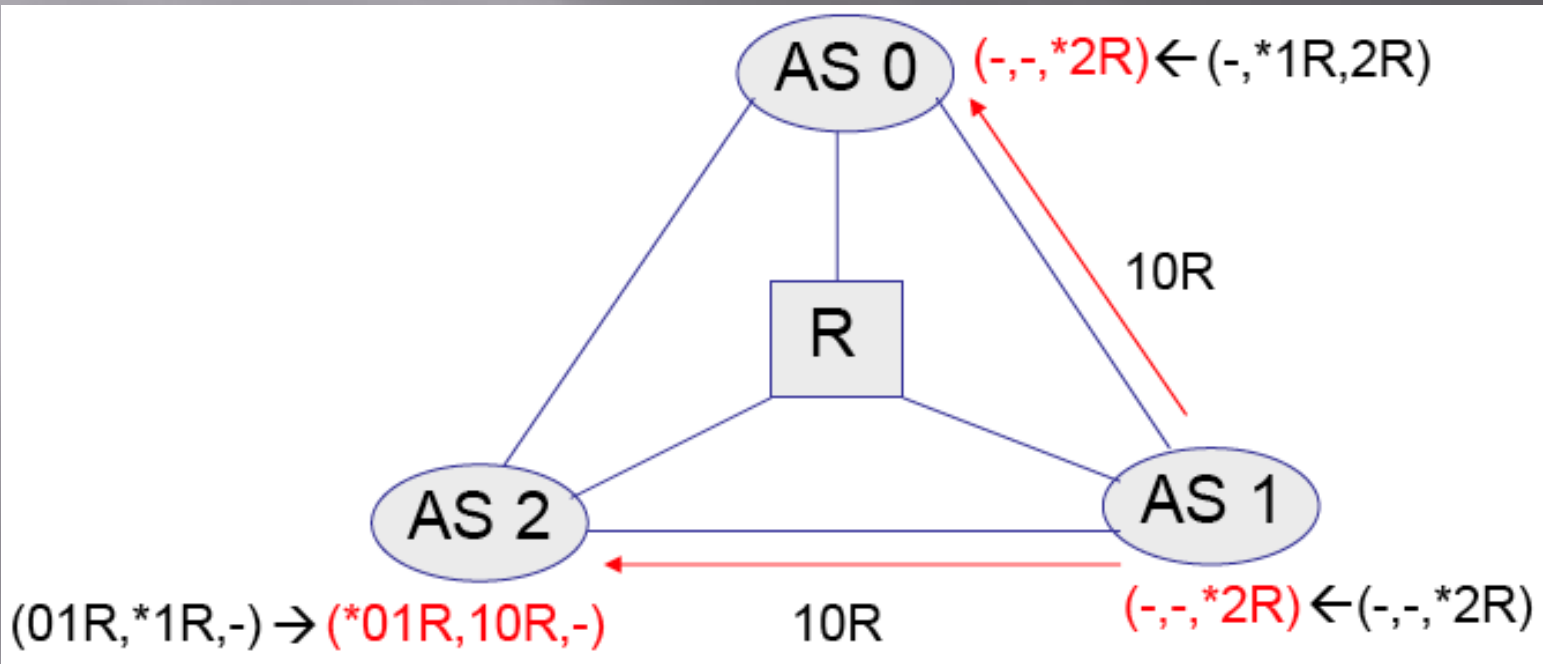
Why “long” failovers?

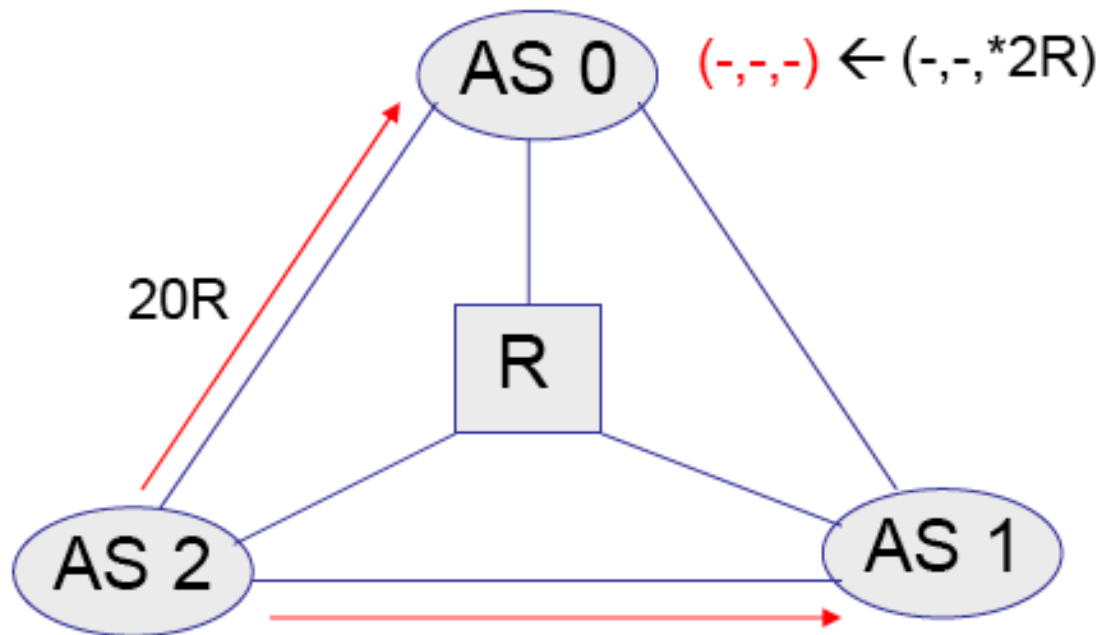
- **Route oscillation**
 - ◆ If policy just uses AS_PATH length, then looks like a DV protocol; can still oscillate
 - ◆ Loop prevention doesn't prevent this
 - ◆ Can explore every possible path through network → (n-1)! Combinations
- **Timers**
 - ◆ 30MinAdver timer makes router wait between updates
 - ◆ Adds artificial “rounds” to propagation speed
- **Loop detection**
 - ◆ Waits to send routes that have loops in them (prevailing BGP implementation only does receiver based loop detection)
- **Typical Internet failover times**
 - ◆ New/shorter link → 60 seconds: simple replacement at nodes
 - ◆ Down link → 180 seconds: search of possible options
 - ◆ Longer link → 120 seconds: replacement or search based on length

Example BGP Oscillation





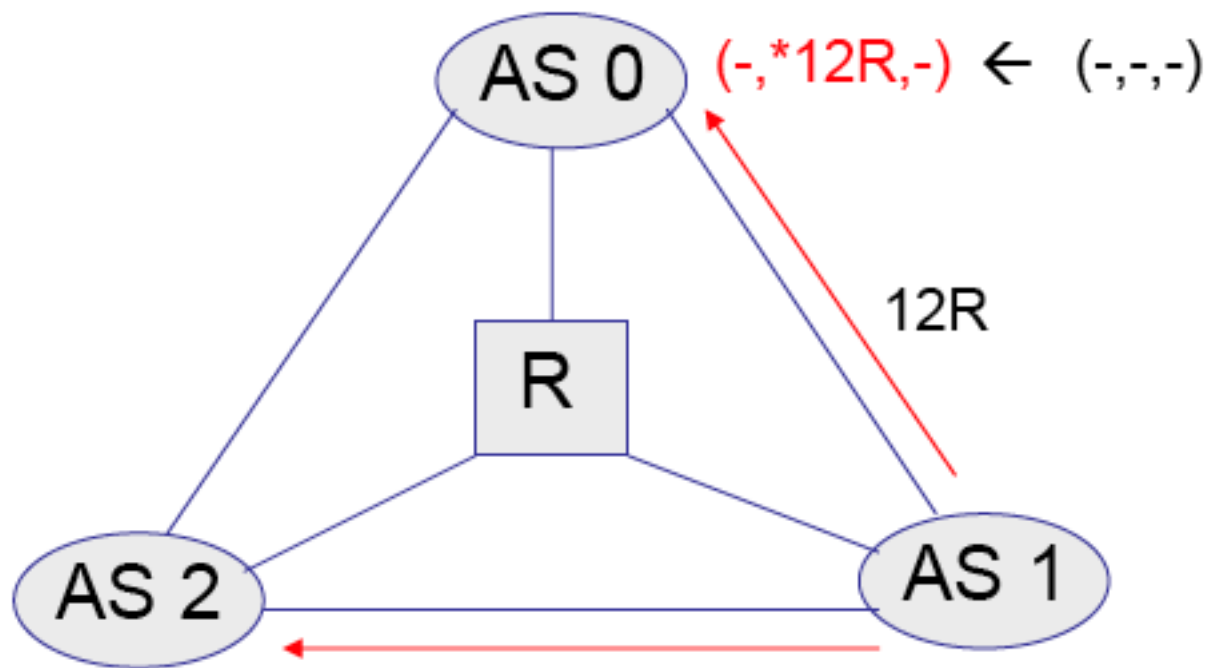




$(*01R, 10R, -) \rightarrow (*01R, 10R, -)$

20R

$(-, -, *20R) \leftarrow (-, -, *2R)$



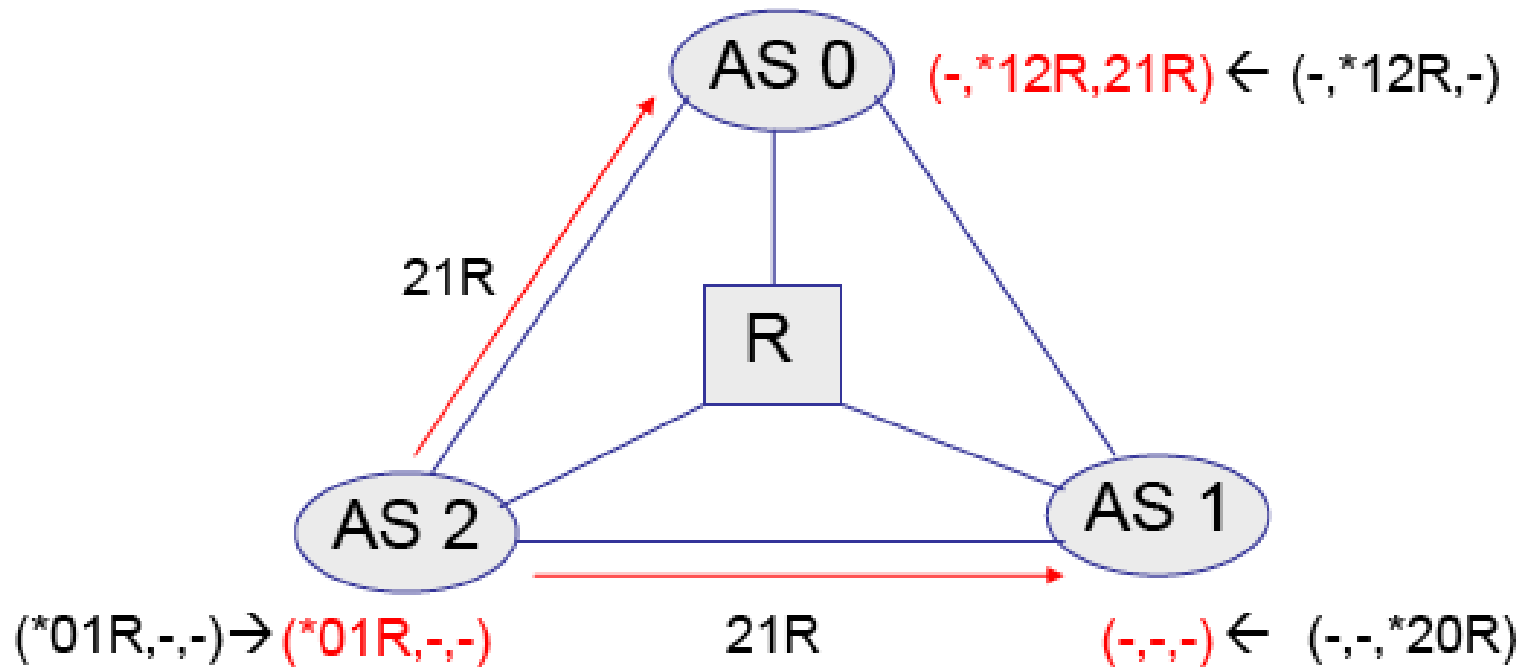
$(*01R, 10R, -) \rightarrow (*01R, -, -)$

12R

$(-, -, *20R) \leftarrow (-, -, *20R)$

$(-, *12R, -) \leftarrow (-, -, -)$

Example Oscillation



BGP impact on flow

- ▣ Slides by Sharad Agarwal

Traffic Collection

- ▣ Packet-level data
 - ▣ Ingress, OC-12
 - ▣ TCP/IP header
 - ▣ GPS synched timestamp
- Analyzed 6 traces
 - ▣ 2 PoPs, 3 days, 6 different ASes